

# Enumeration Schemes for Permutations Avoiding Barred Patterns

Lara Pudwell  
Valparaiso University

Joint Mathematics Meetings  
Washington, DC  
January 5, 2009

# Reduction

- Given a string of numbers  $q = q_1 \cdots q_n$ , the **reduction** of  $q$  is the string obtained by replacing the  $i^{\text{th}}$  smallest number of  $q$  with  $i$ .

# Reduction

- Given a string of numbers  $q = q_1 \cdots q_n$ , the **reduction** of  $q$  is the string obtained by replacing the  $i^{\text{th}}$  smallest number of  $q$  with  $i$ .
- For example, the reduction of 26795 is

# Reduction

- Given a string of numbers  $q = q_1 \cdots q_n$ , the **reduction** of  $q$  is the string obtained by replacing the  $i^{\text{th}}$  smallest number of  $q$  with  $i$ .
- For example, the reduction of 26795 is 1●●●●.

# Reduction

- Given a string of numbers  $q = q_1 \cdots q_n$ , the **reduction** of  $q$  is the string obtained by replacing the  $i^{\text{th}}$  smallest number of  $q$  with  $i$ .
- For example, the reduction of 26795 is 1●●●2.

# Reduction

- Given a string of numbers  $q = q_1 \cdots q_n$ , the **reduction** of  $q$  is the string obtained by replacing the  $i^{\text{th}}$  smallest number of  $q$  with  $i$ .
- For example, the reduction of 26795 is 13●●2.

# Reduction

- Given a string of numbers  $q = q_1 \cdots q_n$ , the **reduction** of  $q$  is the string obtained by replacing the  $i^{\text{th}}$  smallest number of  $q$  with  $i$ .
- For example, the reduction of 26795 is 134●2.

# Reduction

- Given a string of numbers  $q = q_1 \cdots q_n$ , the **reduction** of  $q$  is the string obtained by replacing the  $i^{\text{th}}$  smallest number of  $q$  with  $i$ .
- For example, the reduction of 26795 is 13452.



## Pattern Avoidance/Containment

Given permutations  $p = p_1 \cdots p_n$  and  $q = q_1 \cdots q_m$ ,

- $p$  **contains**  $q$  as a pattern if there is  $1 \leq i_1 < \cdots < i_m \leq n$  so that  $p_{i_1} \cdots p_{i_m}$  reduces to  $q$ ;
- otherwise  $p$  **avoids**  $q$ .

For example,

- 4576213 contains 312 (4**5**762**13**).
- 4576213 avoids 1234.

Notation:  $S_n(q) = \{\pi \in S_n \mid \pi \text{ avoids } q\}$ .

For example,  $|S_n(12)| = 1$  for  $n \geq 0$ .

## Barred Patterns

- Consider *pairs* of permutations  $q_*$ ,  $q^*$  such that  $q_*$  is contained in  $q^*$ .  
e.g.  $q_* = 123$  and  $q^* = 15342$
- Choose one instance of  $q_*$  in  $q^*$ .  
e.g.  $q_* = 123$  and  $q^* = 15342$
- Write  $q$  by taking the letters of  $q^*$  and putting a bar over letters *not* in  $q_*$ .  
For  $q_* = 123$  and  $q^* = 15342$ , we have  $q = 1\bar{5}34\bar{2}$ .

## Barred Pattern Avoidance/Containment

$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

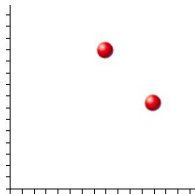
$$|S_n(\overline{132})| = (n-1)!$$

## Barred Pattern Avoidance/Containment

$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$

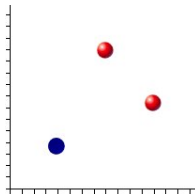


## Barred Pattern Avoidance/Containment

$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$

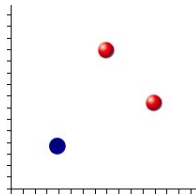


## Barred Pattern Avoidance/Containment

$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$



$$= |\{\pi \in S_n \mid \pi_1 = 1\}|$$

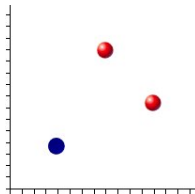
## Barred Pattern Avoidance/Containment

$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$

$$|S_n(\overline{1423})| = B_n \text{ (Bell numbers)}$$



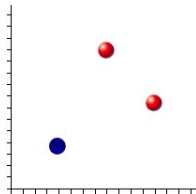
$$= |\{\pi \in S_n \mid \pi_1 = 1\}|$$

## Barred Pattern Avoidance/Containment

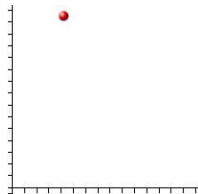
$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$



$$|S_n(\overline{1423})| = B_n \text{ (Bell numbers)}$$



$$= |\{\pi \in S_n | \pi_1 = 1\}|$$

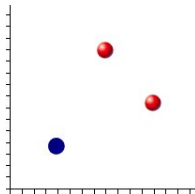


## Barred Pattern Avoidance/Containment

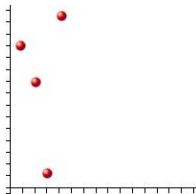
$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$



$$|S_n(\overline{1423})| = B_n \text{ (Bell numbers)}$$



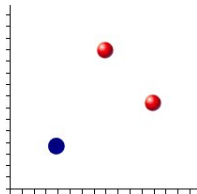
$$= |\{\pi \in S_n | \pi_1 = 1\}|$$

## Barred Pattern Avoidance/Containment

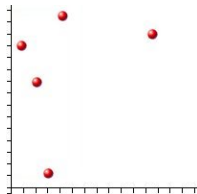
$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$



$$|S_n(\overline{1423})| = B_n \text{ (Bell numbers)}$$



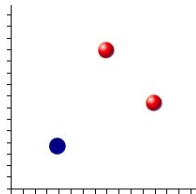
$$= |\{\pi \in S_n \mid \pi_1 = 1\}|$$

## Barred Pattern Avoidance/Containment

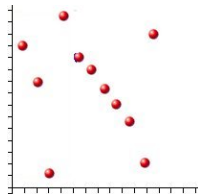
$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$



$$|S_n(\overline{1423})| = B_n \text{ (Bell numbers)}$$



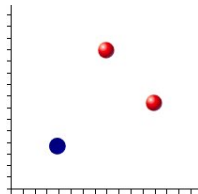
$$= |\{\pi \in S_n | \pi_1 = 1\}|$$

## Barred Pattern Avoidance/Containment

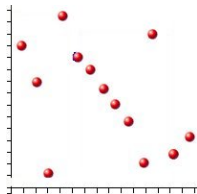
$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$



$$|S_n(\overline{1423})| = B_n \text{ (Bell numbers)}$$



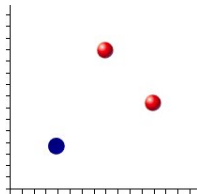
$$= |\{\pi \in S_n \mid \pi_1 = 1\}|$$

## Barred Pattern Avoidance/Containment

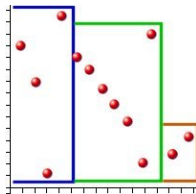
$p$  avoids  $q$  if every instance of  $q_*$  in  $p$  is part of an instance of  $q^*$  in  $p$ .

Some nice examples include:

$$|S_n(\overline{132})| = (n-1)!$$



$$|S_n(\overline{1423})| = B_n \text{ (Bell numbers)}$$



$$= |\{\pi \in S_n \mid \pi_1 = 1\}|$$

## Previous work with Particular Barred Patterns

- West, 1990:  
A permutation is **2-stack sortable**  
if and only if it avoids 2341 and  $3\bar{5}241$ .
- Bousquet-Melou and Butler, 2006:  
A permutation is **forest-like**  
if and only if it avoids 1324 and  $21\bar{3}54$ .
- Claesson, Dukes, and Kitaev, 2008:  
 **$(2 + 2)$ -free posets**  
are in bijection with permutations which avoid  $3\bar{1}52\bar{4}$ .
- Burstein and Lankham, 2006:  
A permutation is a **reverse patience word**  
if and only if it avoids  $3-\bar{1}-42$ .

## Patterns of length $\leq 5$

Length	Bars	No. Sequences	Possible Sequences
2	0	1	1
3	1	2	1, $(n-1)!$
4	2	2	1, $(n-2)!$
5	3	2	1, $(n-3)!$
3	0	1	Catalan
4	1	4	Catalan, Bell, A051295, & A137533
5	2	17	A110447, A117106, & 15 more
4	0	3	A005802, A061552, A022558
5	1	13	A006789, A047970, A098569, A122993, & 9 more

## Patterns of length $\leq 5$

Length	Bars	No. Sequences	Possible Sequences
2	0	1	1
3	1	2	1, $(n-1)!$
4	2	2	1, $(n-2)!$
5	3	2	1, $(n-3)!$
3	0	1	Catalan
4	1	4	Catalan, Bell, A051295, & A137533
5	2	17	A110447, A117106, & <b>15</b> more
4	0	3	A005802, A061552, A022558
5	1	13	A006789, A047970, A098569, A122993, & <b>9</b> more

There are at least 24 new sequences obtained by counting  $S_n(q)$ , where  $q$  is a barred pattern of length 5.



# Goals

## Main Goal

Automate a method to count  $S_n(q)$ .

The method of enumeration schemes is a divide-and-conquer approach to finding recurrences that has been applied to pattern-avoiding permutations (Vatter, Zeilberger), and to pattern-avoiding words (P.).

## Notation

$$S_n(q; p_1 \cdots p_l) := \left\{ \pi \in S_n \mid \begin{array}{l} \pi \text{ avoids } q \\ \pi \text{ has prefix } p_1 \cdots p_l \end{array} \right\}$$

$$S_n\left(q; \begin{array}{c} p_1 \cdots p_l \\ i_1 \cdots i_l \end{array}\right) := \left\{ \pi \in S_n \mid \begin{array}{l} \pi \text{ avoids } q \\ \pi \text{ has prefix } p_1 \cdots p_l \\ \pi = i_1 \cdots i_l \pi_{l+1} \cdots \pi_n \end{array} \right\}$$

For example,

$$S_3(132; 12) = \{123, 231\}$$

$$S_3\left(132; \begin{array}{c} 12 \\ 23 \end{array}\right) = \{231\}$$

## Divide

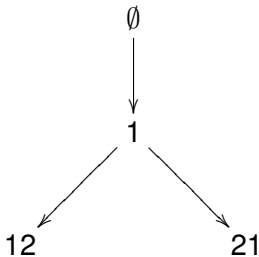
## Notation

$$S_n(q; p_1 \cdots p_l) := \left\{ \pi \in S_n \mid \begin{array}{l} \pi \text{ avoids } q \\ \pi \text{ has prefix } p_1 \cdots p_l \end{array} \right\}$$

$$S_n\left(q; \begin{array}{c} p_1 \cdots p_l \\ i_1 \cdots i_l \end{array}\right) := \left\{ \pi \in S_n \mid \begin{array}{l} \pi \text{ avoids } q \\ \pi \text{ has prefix } p_1 \cdots p_l \\ \pi = i_1 \cdots i_l \pi_{l+1} \cdots \pi_n \end{array} \right\}$$

For any pattern  $q$ , we  
have

$$\begin{aligned} S_n(q) &= S_n(q; 1) \\ &= S_n(q; 12) \cup S_n(q; 21), \\ &\text{etc.} \end{aligned}$$



## Reversibly Deletable

### Reversibly Deletable

Given  $q$  and  $p$ , a position  $r$  is said to be *reversibly deletable* if

- ① deleting  $p_r$  from a  $q$ -avoiding permutation beginning with  $p_1 \cdots p_l$  always produces a  $q$ -avoiding permutation and
- ② inserting  $p_r$  into a  $q$ -avoiding permutation beginning with  $p_1 \cdots p_{r-1} p_{r+1} \cdots p_l$  always produces a  $q$ -avoiding permutation.

If  $r$  is reversibly deletable, we have a natural bijection between  $S_n(q; p_1 \cdots p_r \cdots p_l)$  and  $S_{n-1}(q; p_1 \cdots p_{r-1} p_{r+1} \cdots p_l)$  so that:

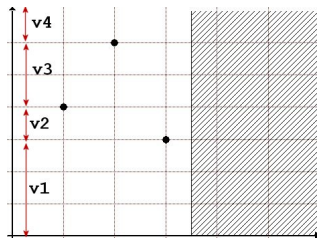
$$|S_n(q; p_1 \cdots p_r \cdots p_l)| = |S_{n-1}(q; p_1 \cdots p_{r-1} p_{r+1} \cdots p_l)|$$

# Spacing Vectors

## Spacing Vectors

Given  $q$  and  $p$  (of length  $l$ ) let  $v$  be a vector in  $\mathbb{N}^{l+1}$ . Then,  $S_n(q; p; v)$  denotes the set of permutations of length  $n$ , avoiding  $q$ , beginning with prefix  $p$  with exactly  $v_1$  letters smaller than "1",  $v_j$  letters greater than " $j-1$ " and smaller than " $j$ ", and exactly  $v_{l+1}$  letters greater than " $l$ ".

For example,

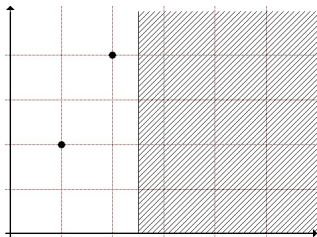


## Spacing Vectors

## Spacing Vectors

Given  $q$  and  $p$  (of length  $l$ ) let  $v$  be a vector in  $\mathbb{N}^{l+1}$ . Then,  $S_n(q; p; v)$  denotes the set of permutations of length  $n$ , avoiding  $q$ , beginning with prefix  $p$  with exactly  $v_1$  letters smaller than “1”,  $v_j$  letters greater than “ $j-1$ ” and smaller than “ $j$ ”, and exactly  $v_{l+1}$  letters greater than “ $l$ ”.

For example,



$$S_n(132; 12; \langle 0, 1, 0 \rangle) = \{ \}.$$

# Gap Vectors

## Gap Vectors

A spacing vector  $v$  is a *gap vector* for  $[q, p]$  if there are no permutations avoiding  $q$  with prefix  $p$  and spacing vector  $\geq v$  (componentwise).

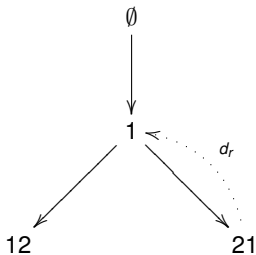
Gap vectors give a condition for which choices of  $i_1, \dots, i_l$  yield

$$\left| S_n \left( q; \begin{matrix} p_1 \cdots p_r \cdots p_l \\ i_1 \cdots i_r \cdots i_l \end{matrix} \right) \right| = 0.$$

# Conquer

## Objectives

- Given  $q$  and  $p$  find  $r$  such that  
 $|S_n(q; p_1 \cdots p_r \cdots p_l)| = |S_{n-1}(q; p_1 \cdots p_{r-1} p_{r+1} \cdots p_l)|$



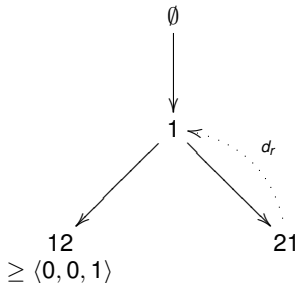
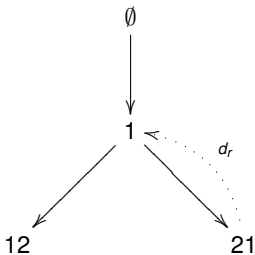


## Conquer

## Objectives

- Given  $q$  and  $p$  find  $r$  such that  
 $|S_n(q; p_1 \cdots p_r \cdots p_l)| = |S_{n-1}(q; p_1 \cdots p_{r-1} p_{r+1} \cdots p_l)|$
- Given  $q$  and  $p$ , find  $i_1, \dots, i_l$  such that

$$\left| S_n \left( q; \begin{matrix} p_1 \cdots p_r \cdots p_l \\ i_1 \cdots i_r \cdots i_l \end{matrix} \right) \right| = 0$$

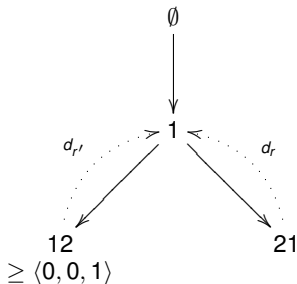
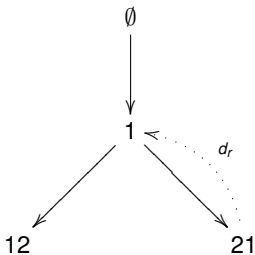


## Conquer

## Objectives

- Given  $q$  and  $p$  find  $r$  such that  
 $|S_n(q; p_1 \cdots p_r \cdots p_l)| = |S_{n-1}(q; p_1 \cdots p_{r-1} p_{r+1} \cdots p_l)|$
- Given  $q$  and  $p$ , find  $i_1, \dots, i_l$  such that

$$\left| S_n \left( q; \begin{matrix} p_1 \cdots p_r \cdots p_l \\ i_1 \cdots i_r \cdots i_l \end{matrix} \right) \right| = 0$$



## Enumeration Scheme

An enumeration scheme is a set of triples  $[p_i, R_i, G_i]$  such that for each triple

- $p_i$  is a reduced prefix of length  $n$
- $R_i$  is a subset of  $\{1, \dots, n\}$
- $G_i$  is a set of vectors of length  $n + 1$  and
- either  $R_i$  is non-empty or all refinements of  $p_i$  are also in the scheme.

## Enumeration Scheme

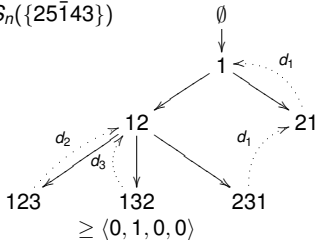
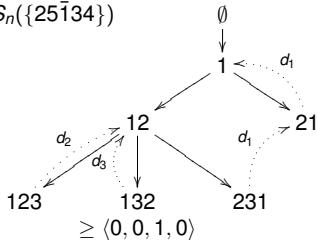
An enumeration scheme is a set of triples  $[p_i, R_i, G_i]$  such that for each triple

- $p_i$  is a reduced prefix of length  $n$  (**prefix**)
- $R_i$  is a subset of  $\{1, \dots, n\}$  (**reversibly deletable positions**)
- $G_i$  is a set of vectors of length  $n + 1$  (**gap vectors**)  
and
- either  $R_i$  is non-empty or all refinements of  $p_i$  are also in the scheme.

# Success Rate

Pattern Lengths	Success Rate	Pattern Lengths	Success Rate
[2,1]	1/1 (100%)	[3,0],[3,0],[3,1]	43/45 (95.6%)
[2,1],[2,0]	2/2 (100%)	[3,0],[3,0],[3,2]	45/45 (100%)
[2,1],[2,1]	2/2 (100%)	[3,0],[3,1],[3,1]	135/138 (97.8%)
		[3,0],[3,1],[3,2]	280/280 (100%)
[3,1]	4/4 (100%)	[3,0],[3,2],[3,2]	138/138 (100%)
[3,2]	4/4 (100%)	[3,1],[3,1],[3,1]	115/118 (97.5%)
[3,0],[3,1]	18/20 (90%)	[3,1],[3,1],[3,2]	378/378 (100%)
[3,0],[3,2]	20/20 (100%)	[3,1],[3,2],[3,2]	378/378 (100%)
[3,1],[3,1]	27/28 (96.4%)	[3,2],[3,2],[3,2]	118/118 (100%)
[3,1],[3,2]	50/50 (100%)		
[3,2],[3,2]	28/28 (100%)	[4,1]	12/16 (75%)
		[4,2]	25/26 (96.2%)
[3,1],[4,0]	59/71 (83.1%)	[4,3]	16/16 (100%)
[3,1],[4,1]	229/240 (95.4%)		
[3,1],[4,2]	355/364 (97.5%)	[5,1]	15/89 (16.9%)
[3,0],[4,1]	84/88 (95.5%)	[5,2]	(in progress)
[3,0],[4,2]	133/136 (97.8%)		
[4,0],[5,1]	(in progress)		

## New Results: Length 5 with 1 Bar

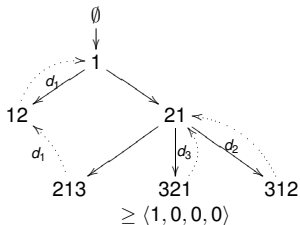
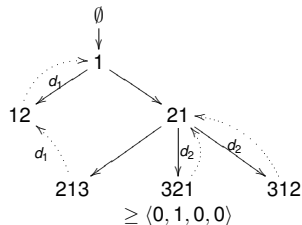
 $S_n(\{25\bar{1}43\})$  $S_n(\{25\bar{1}34\})$ 

both give the sequence

1, 2, 6, 23, 104, 532, 3004, 18426, 121393, 851810, 6325151,  
49448313, 405298482, 3470885747, 30965656442

for  $1 \leq n \leq 15$ .

## New Results: Length 5 with 1 Bar

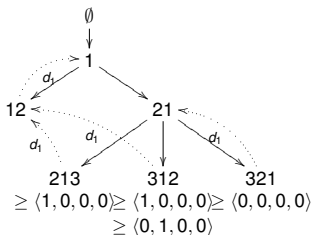
 $S_n(\{43\bar{5}21\})$ 

 $S_n(\{43\bar{5}12\})$ 


both also give the sequence

1, 2, 6, 23, 104, 532, 3004, 18426, 121393, 851810, 6325151,  
49448313, 405298482, 3470885747, 30965656442

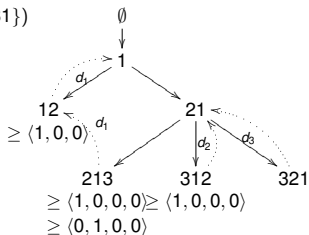
for  $1 \leq n \leq 15$ .

## New Results: Length 5 with 2 Bars

 $S_n(\{\overline{51\bar{2}43}\})$ 

gives the new sequence

1, 2, 5, 14, 43, 143, 511, 1950, 7903,  
33848, 152529, 720466, 3555715,  
18285538, 97752779

for  $1 \leq n \leq 15$ . $S_n(\{\overline{5\bar{4}231}\})$ 

gives the new sequence

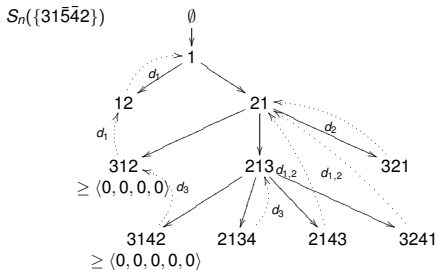
1, 2, 5, 14, 43, 146, 561, 2518, 13563,  
88354, 686137, 6191526, 63330147,  
720314930, 8985750097

for  $1 \leq n \leq 15$ .

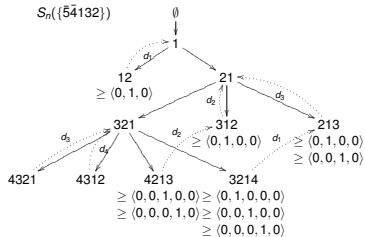


Examples

# New Results: Length 5 with 2 Bars



gives the new sequence  
 1, 1, 2, 5, 14, 43, 144, 522, 2030,  
 8398, 36714, 168793, 813112,  
 4091735, 21451972, 116891160  
 for  $1 \leq n \leq 15$ .



gives the new sequence  
 1, 1, 2, 5, 14, 43, 147, 575, 2648,  
 14617, 96696, 754585, 6794015,  
 69116493, 781266266, 9688636317  
 for  $1 \leq n \leq 15$ .

# Summary

- The method of enumeration schemes confirms many known results for barred patterns and generates new results for  $S_n(25\bar{1}43)$ ,  $S_n(25\bar{1}34)$ ,  $S_n(43\bar{5}21)$ ,  $S_n(43\bar{5}12)$ ,  $S_n(5\bar{1}\bar{2}43)$ ,  $S_n(\bar{5}4\bar{2}31)$ ,  $S_n(31\bar{5}\bar{4}2)$ ,  $S_n(\bar{5}\bar{4}132)$ .
- It remains to find other ways to count permutations avoiding barred patterns.
- There are at least 24 new sequences obtained by counting  $S_n(q)$ , where  $q$  is a barred pattern of length 5.

# Summary

- The method of enumeration schemes confirms many known results for barred patterns and generates new results for  $S_n(25\bar{1}43)$ ,  $S_n(25\bar{1}34)$ ,  $S_n(43\bar{5}21)$ ,  $S_n(43\bar{5}12)$ ,  $S_n(5\bar{1}\bar{2}43)$ ,  $S_n(\bar{5}4\bar{2}31)$ ,  $S_n(31\bar{5}\bar{4}2)$ ,  $S_n(\bar{5}4\bar{1}32)$ .
- It remains to find other ways to count permutations avoiding barred patterns.
- There are at least **19** new sequences obtained by counting  $S_n(q)$ , where  $q$  is a barred pattern of length 5.